

Incidence geometry from an algebraic graph theory point of view

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Outline

- Incidence geometries: examples, problems,...
- Graph theory
- Algebraic techniques in graph theory
and some applications in incidence geometry.

- We study *incidence geometries* (finite or infinite).

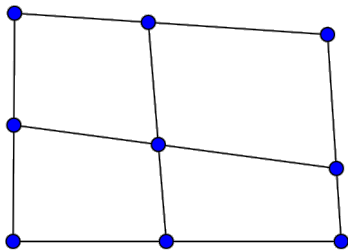
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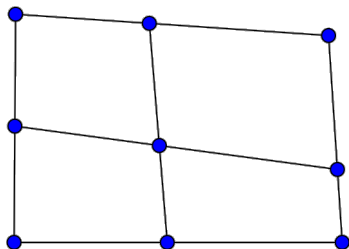
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- Many of these incidence structures are *spherical buildings* (Tits) (related to Coxeter groups)

Example: grids (rank=2)



P_1 : 9 points, P_2 : 6 lines.

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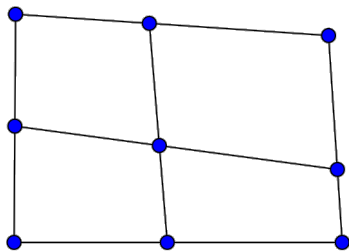


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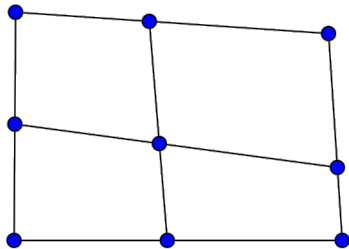


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- Every point is on 2 lines.
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- 2 distinct points are either on 1 common line (collinear) or none.
- If $p \not\in \ell$ then there is a unique $q \in \ell$ collinear with p .

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What is the maximum size for partial spreads of lines?

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- A subspace is *isotropic* if f vanishes on it.
- Let P_i be set of isotropic i -dimensional subspaces.
- Incidence I is just symmetrized inclusion.
- This incidence structure is a *classical polar space*, with rank maximal dimension of isotropic subspaces.

Graphs

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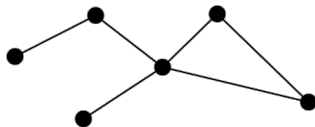
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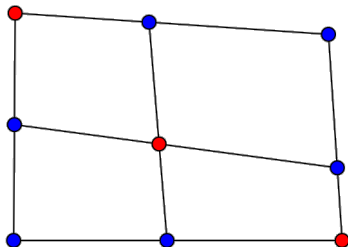
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Cliques and cocliques in graphs

- Clique: subset of vertices, every two elements adjacent.
- Coclique: subset of vertices, no two elements adjacent.

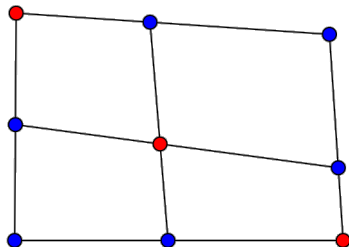
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Vertices: 9 points. Adjacency: collinearity.



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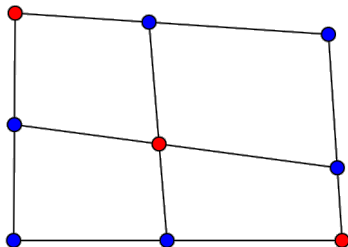
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- Collinearity graph is regular with valency 4.
- The **red points** are a coclique of size 3 (and we can't do better).

Graphs from projective space $PG(n, q)$

Grassmann graph $J_q(n+1, 2)$ in $V(n+1, q)$:

- 2-dimensional subspaces as vertices
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Finding partial spreads of lines in $\text{PG}(n, q) =$

Finding cliques in $J_q(n+1, 2)$.

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Dual polar graph:

- vertices: isotropic subspaces of maximal dimension d
- x and y are adjacent if $x \cap y$ is $(d - 1)$ -dimensional.

Consider graph Γ with vertex set Ω :

- Adjacency matrix A : $(|\Omega| \times |\Omega|)$ -matrix:

$$(A_i)_{x,y} = \begin{cases} 1 & \text{if } x \text{ and } y \text{ are adjacent} \\ 0 & \text{if } x \text{ and } y \text{ are not adjacent} \end{cases} .$$

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- Eigenvalues of Γ* = the eigenvalues of A .
- A is real-symmetric \implies all eigenvalues are real.
- If Γ is regular with valency k then $A(1, 1, \dots, 1)^T = k(1, 1, \dots, 1)^T$, so k is eigenvalue (and largest eigenvalue).

Characteristic vectors

For any subset of vertices S ,

the *characteristic vector* $\chi_S \in \mathbb{R}^\Omega$ is $(0, 1, 0, 1, 1, \dots)^T$

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Orthogonal decomposition into eigenspaces

We can write:

$$I = E_1 + \dots + E_d,$$

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Just $\langle E_i\chi_S, E_i\chi_S \rangle \geq 0$ might already be powerful!

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- This yields new (tight) bound for one specific type of polar space!

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- Nice things happen in case of equality!

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- Joint work with V. Pepe and L. Storme: using properties in case of equality to characterize Erdős-Ko-Rado-sets of maximum size.
- Characterization was possible for almost all polar spaces, easy construction is unique optimal construction in most cases.

Thank you for your attention!

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