

Some recent developments in the theory of linear MRD-codes

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Outline

- 1 Linear MRD-codes
- 2 Generalized Gabidulin Codes and linearized polynomials
- 3 Generalized Twisted Gabidulin Codes
- 4 MRD-codes-Maximum Scattered Spaces-Segre Variety

RD-codes

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d(C) minimum distance of C

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$Aut(\mathcal{C})$ Automorphism group of \mathcal{C}

$$Aut(\mathcal{C}) = \{(A, B, \sigma) \in GL(m, q) \times GL(n, q) \times Aut(\mathbb{F}_q) : AC^\sigma B = \mathcal{C}\}$$

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$$L(\mathcal{C})^* \times R(\mathcal{C})^* \times \{id\} \subseteq Aut(\mathcal{C})$$

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If $m \leq n$ then $L(\mathcal{C})$ is a finite field. Hence $|L(\mathcal{C})| \leq q^m$.
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$\mathcal{C}_{\mathcal{G}_{k,s}}(U)^T$	Generalised Gabidulin MRD-code	$[m \times n, kn, n - k + 1]_{\mathbb{F}_q}$

Linearized polynomials

$$\mathbb{F}_q^{n \times n}$$

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$$\mathcal{L}_{n,q} = \{f(x) = a_0x + a_1x^q + \dots + a_{n-1}x^{q^{(n-1)}} : a_0, a_1, \dots, a_{n-1} \in \mathbb{F}_{q^n}\}$$

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\circ \rightarrow composition modulo $x^{q^n} - x$

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Finite presemifields

J. De La Cruz, Kiermaier, Wassermann, Willem: Algebraic structures of MRD codes, *Adv. Math. Commun.* (2016)

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$\mathcal{G}_{1,s}$ \rightarrow $\mathcal{S}_{\mathcal{G}_{1,s}}$ isotopic to \mathbb{F}_{q^n}

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Generalized Twisted Gabidulin Codes

Sheekey, 2016

$$\begin{aligned}\mathcal{H}_{k,s}(\eta, h) = \{f(x) = a_0x + a_1x^{q^s} + \cdots + a_{k-1}x^{q^{s(k-1)}} + \eta a_0^{q^h}x^{q^{sk}} : a_i \in \mathbb{F}_{q^n}\} \\ n, k, s \in \mathbb{Z}^+, \quad k < n, \quad \gcd(n, s) = 1, \quad N_q(\eta) \neq (-1)^{nk}\end{aligned}$$

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$$\mathcal{H}_2(\eta, 1)$$

K. Otal, F. Özbudak: Explicit Construction of Some Non-Gabidulin Linear Maximum Rank Distance Codes, *Adv. Math. Commun.* (2016)

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Sheekey 2016/ Lunardon-Trombetti-Zhou, ArXiv (2015)

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Horlemann-Trautmann, Marshall: New Criteria for MRD and Gabidulin Codes and some Rank-Metric Code Constructions, arXiv:1507.08641v3 (2016)

MRD-codes and maximum scattered spaces

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$$\Phi_f(x) = \frac{f(x)}{x} \quad \longrightarrow \quad |Im(\Phi_f)| = q^{n-1} + q^{n-2} + \cdots + q + 1$$

MRD-codes and maximum scattered spaces

Sheekey, AMC (2016)

$$\mathcal{C} \in \mathcal{M}_L, \quad k = 2$$

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Maximum scattered spaces-Maximum scattered linear sets

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Two-weight linear codes \longleftrightarrow Projective two-intersection sets

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R. Calderbank, W.M. Kantor: The geometry of two-weight codes, *Bull. London Math. Soc.* (1986)

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- $U_1 = \{(x, x^{q^s}) : x \in \mathbb{F}_{q^n}\}, 1 \leq s \leq n-1, \gcd(s, n) = 1.$
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$$k = 2, n = 6, 8 \quad \Rightarrow \quad \mathcal{H}_L \subset \mathcal{M}_L$$

MRD-codes and Maximum Scattered Spaces

Classification results

- \mathcal{C} MRD-code, $[3 \times 3, 6, 2]_{\mathbb{F}_q}$, \mathbb{F}_{q^3} -linear

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For $q = 2$ - **Horlemann-Trautmann, Marshall:** New criteria for MRD and Gabidulin codes and some rank-metric code constructions, arXiv:1507.08641v3 (2015)

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$\mathcal{C} \subset End(\mathbb{F}_{q^n}, \mathbb{F}_q)$ linear **MRD-code**, $[n \times n, kn, n - k + 1]_{\mathbb{F}_q}$



$$P(\mathcal{C}) \Rightarrow P(\mathcal{C})^\Phi = P(\mathcal{C}^T)$$

$$P(\mathcal{C}) \Rightarrow P(\mathcal{C})^\tau = P(\mathcal{C}^\perp)$$

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- Understand the relationship between Sheekey's construction and Lunardon's approach.

Thank you

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