

# Commutative Semifields

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# Origin

- Finite semifields are related to many important geometric and combinatorial constructs.
- Commutative semifields are of particular interest, as historically there have been few known examples.
- The problem is approached using the connection between commutative semifields and certain types of linear sets.
- We detail computational results towards classifying examples with rank 2 and 3 over the left nucleus (R2CS and R3CS).

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# Finite semifields

A **semifield**  $\mathbb{S}$  is a possibly non-associative algebra with an identity and no zero divisors.

The **left nucleus**  $N_\ell(\mathbb{S})$  (right, middle) is the set of elements  $x \in \mathbb{S}$  such that, for all  $y, z \in \mathbb{S}$ ,  $x \circ (y \circ z) = (x \circ y) \circ z$ .

The intersection of the left, right, and middle nuclei is called simply the **nucleus**. The intersection of the nucleus with the commutative center is called the **center**  $Z(\mathbb{S})$ .

The **dimension** of a semifield is the (vector space) dimension over its center; the **rank** of a semifield is its dimension the left nucleus.

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# Semifield spreads

Semifields can be described in terms of **spread sets** of  $q^k \times k$  matrices over  $\mathbb{F}_q$  with nonsingular pairwise differences.

A spread set over  $\mathbb{F}_q$  determines a spread  $\mathcal{S}$  of  $\text{PG}(2k-1, q)$ , giving a translation plane  $A(\mathcal{S})$  of order  $q^k$ .

A spread  $\mathcal{S}$  is a **semifield spread** if the spread set is an additive subgroup of  $\mathcal{M}_k(q)$ , in which case  $A(\mathcal{S})$  can be coordinatized by a semifield of order  $q^k$  whose left nucleus contains  $\mathbb{F}_q$ .

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# Knuth orbit and commutativity

The **Knuth orbit** of  $\mathbb{S}$  is a collection of at most six isotopism classes  $\{[\mathbb{S}], [\mathbb{S}^t], [\mathbb{S}^d], [\mathbb{S}^{td}], [\mathbb{S}^{dt}], [\mathbb{S}^{tdt}]\}$ ,  $t$  and  $d$  denote *transpose* and *dual* operations.

Commutativity of a semifield is **not** invariant under isotopism, a semifield  $\mathbb{S}$  can be considered “**commutative**” if it is isotopic to a commutative semifield, i.e. if  $[\mathbb{S}^d] = [\mathbb{S}]$ .

Then the subspaces of the spread associated with  $\mathbb{S}^{td}$  are totally isotropic with respect to a symplectic polarity of  $\text{PG}(2k - 1, q)$ , i.e.  $\mathbb{S}^{td}$  is a **symplectic** semifield.

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# Linear sets

Under the field reduction map,  
points of  $\text{PG}(t-1, q^n) \rightarrow$  Desarguesian spread of  $\text{PG}(tn-1, q)$ .

An  $\mathbb{F}_q$ -**linear set**  $\mathcal{L}$  of  $\text{PG}(t-1, q^n)$  corresponds to the spread elements intersecting some subspace  $U$  of  $\text{PG}(tn-1, q)$ . The **rank** of  $\mathcal{L}$  is  $\dim(U)$ .

$$\begin{array}{ll} \text{Semifield } \mathbb{S}, |\mathbb{S}| = q^{kn} & \longleftrightarrow \text{Spread set } \mathcal{S} \subseteq \text{PG}(k^2-1, q^n) \\ \mathbb{F}_{q^k} \subseteq \mathbb{N}_\ell(\mathbb{S}), \mathbb{F}_q \subseteq \mathbb{Z}(\mathbb{S}) & \mathbb{F}_q\text{-linear, disjoint from } \mathcal{S}_{k,k}(q^n) \end{array}$$

If  $\mathbb{S}$  is symplectic, (WLOG) the matrices of  $\mathcal{S}$  are symmetric, contained in a  $\left(\frac{k(k+1)}{2} - 1\right)$ -dimensional subspace intersecting  $\mathcal{S}_{k,k}(q^n)$  in a quadratic Veronesean  $\mathcal{V}_k$ .

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# Rank 2 commutative semifields

Easiest place to start is with commutative semifields having rank 2 over their left nucleus (R2CS).

Cohen and Ganley showed that any R2CS of order  $q^{2n}$  with center  $\mathbb{F}_q$ ,  $q$  odd, arises from a pair  $(f, g)$  of  $\mathbb{F}_q$ -linear functions such that  $g^2(t) - 4tf(t)$  is a nonsquare for all  $t \in \mathbb{F}_{q^n}^*$ .

This is equivalent to the existence of a rank  $n$   $\mathbb{F}_q$ -linear set

$$\mathcal{W} = \{(t, f(t), g(t)) : t \in \mathbb{F}_{q^n}^*\}$$

contained in the set of interior points  $\mathcal{I}(\mathcal{C})$  of the conic  $\mathcal{C}$  with equation  $X_2^2 - 4X_0X_1 = 0$  in  $\text{PG}(2, q^n)$ .

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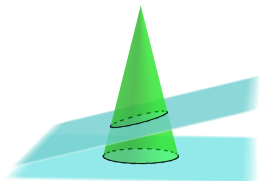
# Semifield flocks and translation ovoids

These functions  $f$  and  $g$  give connections to many other important geometric objects. For example

- The functions  $f$  and  $g$  defining a R2CS determine planes

$$\{\pi_t : tX_0 + f(t)X_1 + g(t)X_2 + X_3 = 0 \mid t \in \mathbb{F}_{q^n}\}$$

in  $\text{PG}(3, q^n)$  forming a **semifield flock of a quadratic cone**.



- A semifield flock of a quadratic cone in  $\text{PG}(3, q^n)$  is equivalent to a **translation ovoid** of the GQ  $\mathcal{Q}(4, q^n)$ .



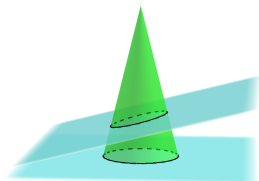
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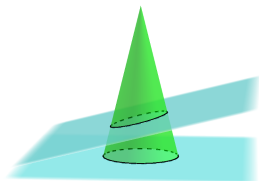
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# Known R2CS examples and bounds

There are no R2CS with  $q$  even. For  $q$  odd, there are very few known examples.

- The Dickson semifields (1906).
- The Cohen-Ganley semifields (1982), which have order  $3^{2n}$  for  $n \geq 2$  and center  $\mathbb{F}_3$ .
- The example found by Penttila and Williams (2000), having order  $3^{10}$  and center  $\mathbb{F}_3$ .

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# Structure of linear sets

If the linear set  $\mathcal{W}$  associated with an R2CS is contained in a line, then it must be a Dickson semifield; to find new semifields we need linear sets that contain an  $\mathbb{F}_q$ -subplane.

Theorem (Ball, Blokhuis, Lavrauw 2003; Lavrauw, 2006)

*If there exists an  $\mathbb{F}_q$ -subplane in  $\text{PG}(2, q^n)$  contained in  $\mathcal{I}(\mathcal{C})$  then there exists an  $\mathbb{F}_q$ -subline contained in  $\ell \cap \mathcal{I}(\mathcal{C})$  with  $\ell$  external to  $\mathcal{C}$ ; such a subline does not exist for*

$$q \geq 4n^2 - 8n + 2,$$

*or for*

$$q > 2n^2 - (4 - 2\sqrt{3})n + (3 - 2\sqrt{3})$$

*when  $q$  is prime.*

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# Outline of the algorithm

This leads us to the following strategy for searching for  $2k$ -dimensional R2CS:

- Determine values of  $q$  for which there actually exists an  $\mathbb{F}_q$ -subline contained in  $\mathcal{I}(\mathcal{C})$  spanning an external line to  $\mathcal{C}$  (for some fixed conic  $\mathcal{C}$ ).
- For these values, find all  $\mathbb{F}_q$ -sublines contained in  $\mathcal{I}(\mathcal{C})$  and determine whether there are two which generate a suitable  $\mathbb{F}_q$ -subplane.
- Use a clique-finding algorithm to determine if subplanes can be combined to give a rank  $k$   $\mathbb{F}_q$ -linear set contained in  $\mathcal{I}(\mathcal{C})$ .

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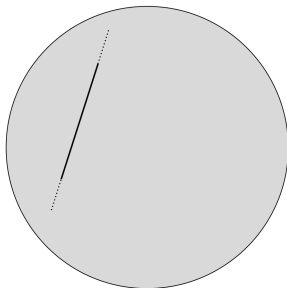
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- Use a clique-finding algorithm to determine if subplanes can be combined to give a rank  $k$   $\mathbb{F}_q$ -linear set contained in  $\mathcal{I}(\mathcal{C})$ .

# Sublines

Want to find  $\mathbb{F}_q$  sublines contained in  $\mathcal{I}(\mathcal{C})$  spanning a line external to  $\mathcal{C}$ , using the group of the conic we can fix an external line  $\ell_e$  and a point  $\mathbf{x}$ , restrict our search for sublines of  $\ell_e$  containing  $\mathbf{x}$ .



# Subplanes and higher dimensional spaces

Looking for  $\mathbb{F}_q$ -subplanes in  $\mathcal{I}(\mathcal{C})$ , again restrict to subplanes containing  $\mathbf{x}$ .

We want **all**  $\mathbb{F}_q$ -sublines on  $\mathbf{x}$ , whether they span a secant or an external line to  $\mathcal{C}$ .

For each pair of sublines, generate points of their subplane one-by-one, check they are in  $\mathcal{I}(\mathcal{C})$ ; if the test fails for a single point reject the pair as incompatible.

For  $\ell$  occurring first in a compatible pair, define a graph on the sublines occurring second with adjacency given by compatibility.

A rank  $k$   $\mathbb{F}_q$ -linear set will correspond to a  $\begin{bmatrix} k-1 \\ 1 \end{bmatrix}_q - 1$  clique in such a graph, so we search for cliques of the appropriate size.

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# Results

Number of sublines on $\mathbf{x} \in \ell_e$			
$q$	$n = 3$	$n = 4$	$n = 5$
3	12	120	1200
5	12	600	15072
7	24	912	52080
9	0	1040	91880
11	0	744	115572
13	0	504	102340
17	-	72	$\geq 1$
19	-	80	$\geq 1$
23	-	0	$\geq 1$
25	-	0	$\geq 1$
27	-	0	$\geq 1$
29	-	0	$\geq 1$
31	-	-	$\geq 1$

# 8-dimensional

An 8-dimensional R2CS is equivalent to a rank 4 linear set contained in  $\mathcal{I}(\mathcal{C})$  for some conic  $\mathcal{C}$  in  $\text{PG}(2, q^4)$ .

While the bound tells us we could have  $q < 30$ ,  $q = 19$  is the largest value for which there is a suitable subline.

There are subplanes only when  $q = 3$ ; have 237 graphs to test, the largest containing 204 vertices. Obtain 174 total cliques of size 12 all giving equivalent rank 4 linear sets (Cohen-Ganley R2CS).

## Theorem

*An 8-dimensional R2CS is either a Dickson semifield, or of Cohen–Ganley type (with center  $\mathbb{F}_3$ ).*

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Searching for suitable  $\mathbb{F}_q$ -linear sublines in  $\text{PG}(2, q^5)$  is more computationally difficult; as  $q$  starts to grow, we cannot search exhaustively.

We can complete the search for  $q = 3$ , finding two nonequivalent rank 5 linear sets (Cohen–Ganley and the Penttinen–Williams examples).

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# Rank 3 commutative semifields

The situation for semifields having rank 3 over the left nucleus is more complicated.

For the case of semifields 6-dimensional over center  $\mathbb{F}_q$ , we need rank 6  $\mathbb{F}_q$ -linear sets in  $\text{PG}(5, q^2)$  disjoint from the secant variety of the quadratic Veronesean  $\mathcal{V}$ .

There are no known examples of 6-dimensional R3CS when  $q$  is even. When  $q$  is odd, we have

- the twisted fields (Albert, 1961);
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# A bound on $q$

Due to a recent result of Marino and Pepe (2016), we do have a bound on  $q$ .

They show that, when  $q^2 > 2 \cdot 3^8$  (so  $q > 114$ ) the list of R3CS from the previous slide is complete.

Furthermore, new examples for  $q < 114$  must correspond to a  $\mathbb{F}_q$ -linear set consisting of  $q^3 + q^2 + q + 1$  lines passing through a common point (a cone with base  $\text{PG}(3, q)$ ).

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# Searching for linear sets

To extend our work on R2CS to the case of R3CS, we adapted our algorithms to the new setting of the quadratic Veronesean.

Searching for a suitable  $\mathbb{F}_3$ -linear set in  $\text{PG}(5, 9)$ , we worked in the quotient space wrt a point  $\mathbf{x}$ . Our task was then to find a rank 4 linear set (subgeometry) contained in a set  $\mathcal{U}$  of “allowed” points.

We build such a set by adding a point at a time, forming a  $\mathbb{F}_q$ -basis for the linear set, and removing orbits of points from  $\mathcal{U}$  as we went.

After a few months(!) of computer time:

## Theorem

*A 6-dimensional rank 3 commutative semifield with center  $\mathbb{F}_3$  is isotopic to either a twisted field, or one of the examples given by BH, LMPT, ZP.*

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# Final remarks

Our techniques worked very well for classifying the 8-dimensional R2CS, but the 10-dimensional classification is much more difficult.

- Can we improve the bound on  $q$  in terms of  $n$ ?
- Can we impose stronger structural requirements on the linear set  $\mathcal{W}$ ?
- Is it even possible to have subplanes contained in  $\mathcal{I}(\mathcal{C})$  when  $q \not\equiv 0 \pmod{3}$ ?

For R3CS, the clique-searching algorithm was completely impractical. We likely need better structural constraints (maybe on  $\langle \mathcal{L} \rangle \cap \mathcal{V}$ ) to proceed.

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