

# SMALL WEIGHT CODE WORDS

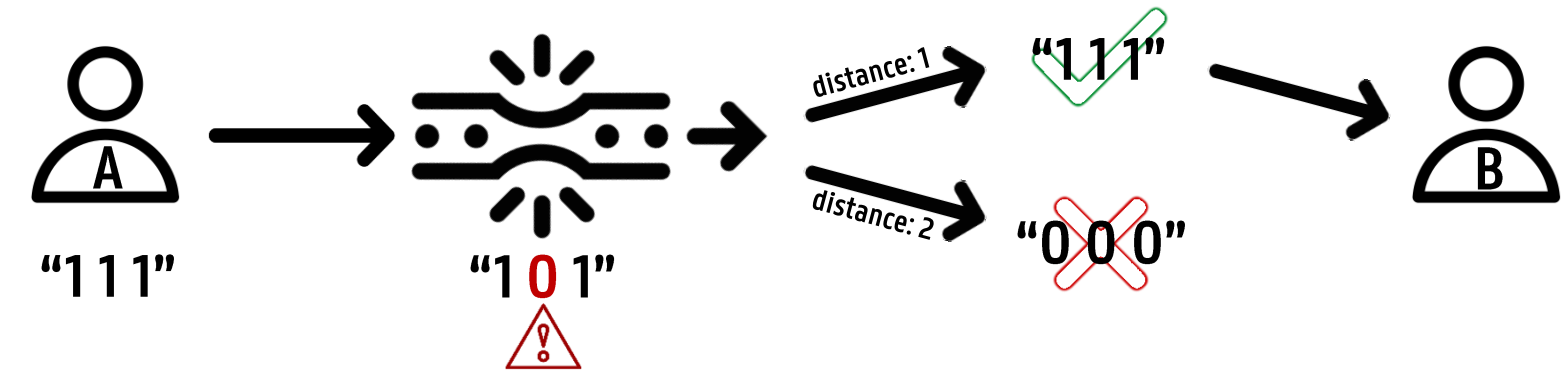
arising from the incidence matrix of points and hyperplanes



## Root of communication

Code words have been used throughout the ages to **transmit a message from point A to point B**. During this transmission, a code word can alter and errors can occur due to external factors (e.g. a hardware defect).

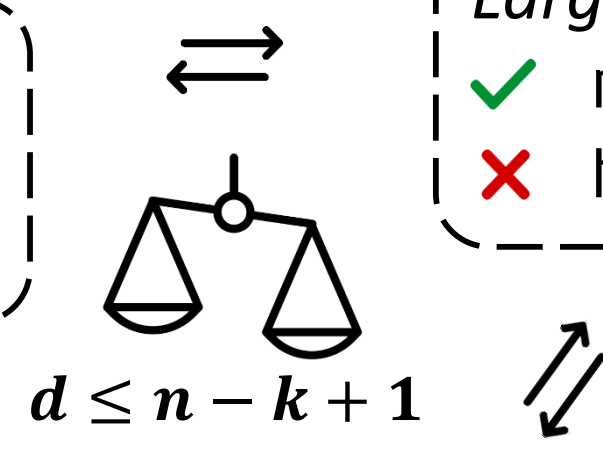
Making a good choice of which code to use and what (de)coding methods to apply, error-correction can prevent the loss of information.



## Linear codes

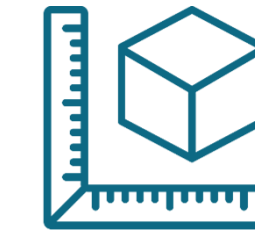
A  $[n, k, d]$ -linear code over the field  $F_q$  is a subspace of the vector space  $V(n, q)$ . Its vectors are called **code words**.

Large length  $n$ :  
 ✓ rich code words  
 ✗ long transmission time



Large dimension  $k$ :  
 ✓ more code words  
 ✗ hard to distinguish

Large minimal distance  $d$ :  
 ✓ error-correcting  
 ✗ large code



## Distance and weight

The (Hamming) **distance** between two code words is measured as the number of different digits arising in both vectors.

The **weight** of a code word is the number of non-zero digits it contains.

The **minimal distance**  $d$  of a code is the smallest distance between two different code words. For linear codes, this is equal to the **minimal weight** of the code.



## Coding theory

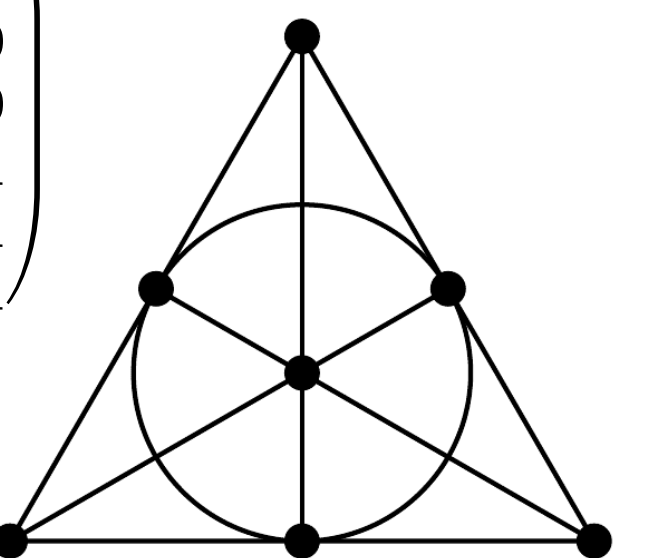
**Coding theory** analyses codes, their properties and usage in various applications. The study of 'good' error-correcting codes, like mentioned above, is an important topic within this field of study.



## The code $C_{n-1}(n, q)$

Numbering the points and hyperplanes of the projective space  $PG(n, q)$ , we can compose the incidence matrix between these two types of spaces.

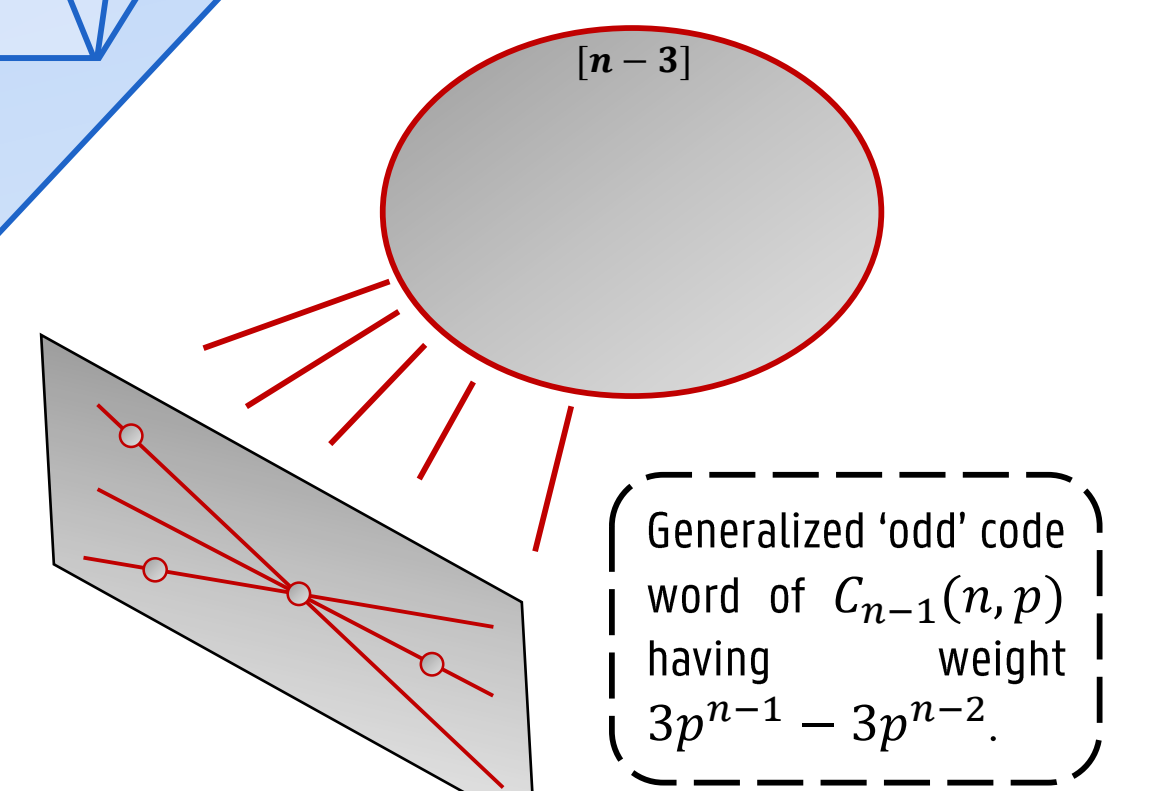
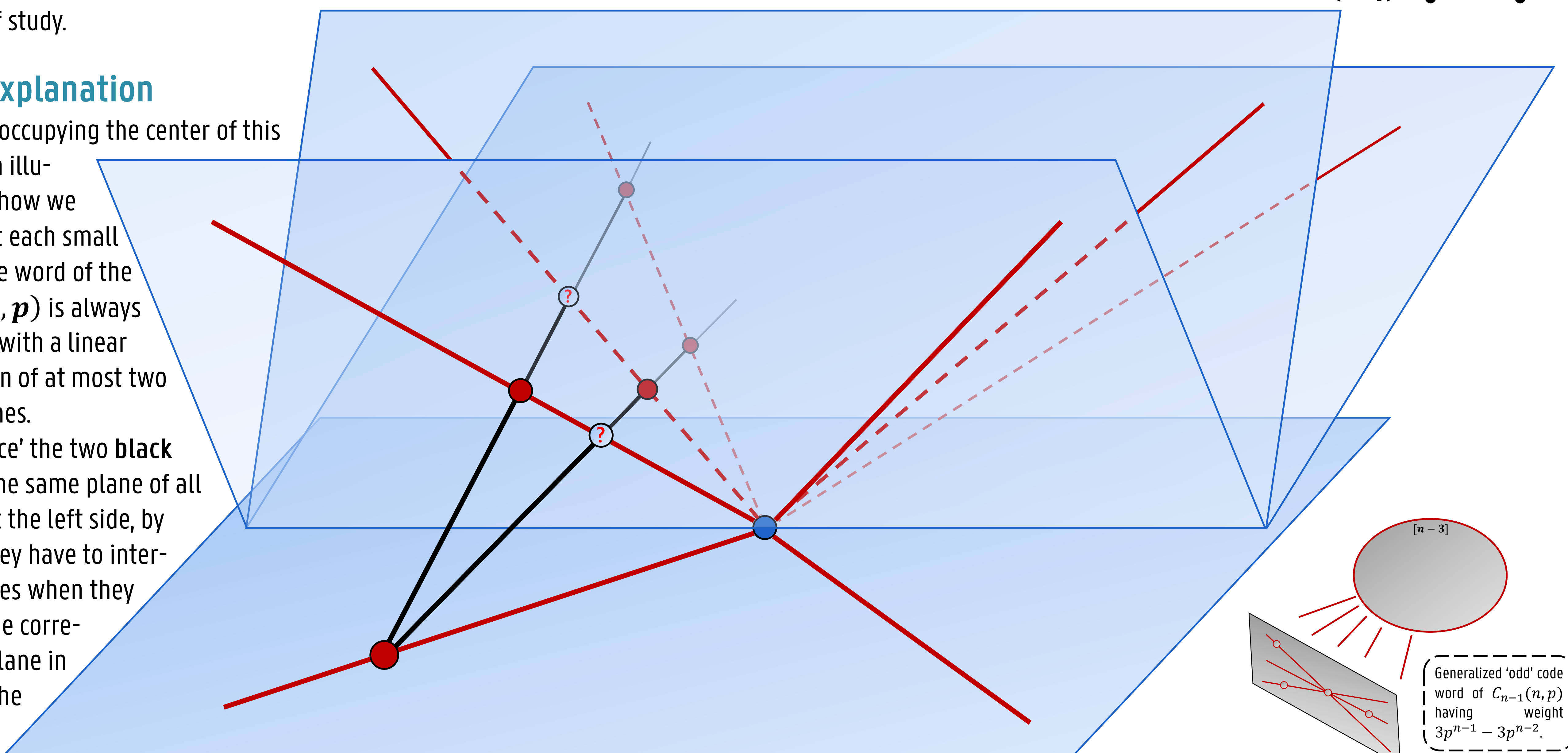
The code  $C_{n-1}(n, q)$  is equal to the span of the rows of this incidence matrix. As such, each code word can be interpreted as a **'linear combination'** of hyperplanes in  $PG(n, q)$ .

$$\begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}$$


## Figure explanation

The **figure** occupying the center of this poster is an illustration of how we proved that each small weight code word of the code  $C_2(3, p)$  is always associated with a linear combination of at most two (hyper)planes.

We can 'force' the two **black** lines into the same plane of all **red** lines at the left side, by knowing they have to intersect **red** lines when they intersect the corresponding plane in a point of the code word.



## Results for the plane

Below are some results concerning  $C_1(2, q)$ . For  $q = p$  prime, all code words of weight at most  $O(3p)$  are associated with a linear combination of at most two lines.

For  $q$  not prime, all code words of weight at most  $O(\frac{1}{2}q^{1.5})$  are associated with a linear combination of at most  $O(\frac{1}{2}\sqrt{q})$  lines.



## Results for arbitrary dimension

It is a well known result that the minimum weight of  $C_{n-1}(n, q)$  equals the number of points in a hyperplane ( $= \theta_{n-1}$ ).

Recently, it was proved that the code words having second smallest weight (namely  $2q^{n-1}$ ) are associated with **the symmetric difference of two different hyperplanes**.



## An 'odd' code word of $C_1(2, p)$

Contrary to our intuition, one can construct a code word of weight  $3p - 3$ , **not associated with a linear combination of three lines**. Generalizing this result, we deduce that small weight code words do not always arise from linear combinations of few hyperplanes, in contrary to speculation.



## The third smallest weight

Our main result states that the **third smallest weight** code words of  $C_{n-1}(n, q)$  have weight  $2q^{n-1} + \theta_{n-2}$ . Such code words are associated with a linear combination (not the difference) of two different hyperplanes.



## The fourth smallest weight?

We have deduced **several lower bounds** for the fourth smallest weight  $w_4$  in the code  $C_{n-1}(n, q)$ , depending on the value  $q$ :

- $w_4 \geq O(\frac{5}{2}q^{n-1})$  when  $q$  is prime.
- $w_4 \geq O(\frac{1}{2}q^{n-1}\sqrt{q})$  otherwise.



## Future work

We will try to **optimise** our lower bounds and **classify** the small weight code words of  $C_{n-1}(n, p)$  as being either linear combinations of hyperplanes or generalizations of the 'odd' code word mentioned above.